

A Delsarte-Style Proof of the Bukh–Cox Bound

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Outline

- 1 The Bukh–Cox Bound
- 2 The Bukh–Cox Lemma
- 3 Bukh–Cox Bound via Linear Programming

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Introduction

- The Bukh–Cox bound addresses the line packing problem: Pack n points in $\mathbb{R}\mathbf{P}^{d-1}$ ($\mathbb{C}\mathbf{P}^{d-1}$) that maximizes the minimum distance.
- Instances of this problem can be traced back to Tammes and F. Tóth.
- Most recent results identify new packings which achieve equality in the Welch Bound.
- Bukh and Cox discovered a different bound, along with a family of packings which achieve said bound.

Definition

Let $X = \{x_i\}_{i \in [n]}$ be a sequence of unit vectors in \mathbb{C}^d . We define the *coherence* of X to be

$$\mu(X) := \max_{i \leq i < j \leq n} |\langle x_i, x_j \rangle|.$$

Essentially, $\mu(X)$ computes the smallest angle between any two vectors in X .

Welch Bound

Theorem (Welch '74)

Let $n > d$ and let $X = \{x_i\}_{i \in [n]}$ be a sequence of unit vectors in \mathbb{C}^d . Then,

$$\mu(X) \geq \sqrt{\frac{n-d}{d(n-1)}},$$

where equality is achieved if and only if X is an equiangular tight frame in \mathbb{C}^d .

Bukh–Cox Bound

Theorem (Bukh, Cox '18)

Let $n > d$ and let $X = \{x_i\}_{i \in [n]}$ be a sequence of unit vectors in \mathbb{C}^d .
Then,

$$\mu(X) \geq \frac{(n-d)^2}{n + (n^2 - nd - n)\sqrt{1 + n - d} - (n-d)^2}.$$

Comparing the bounds

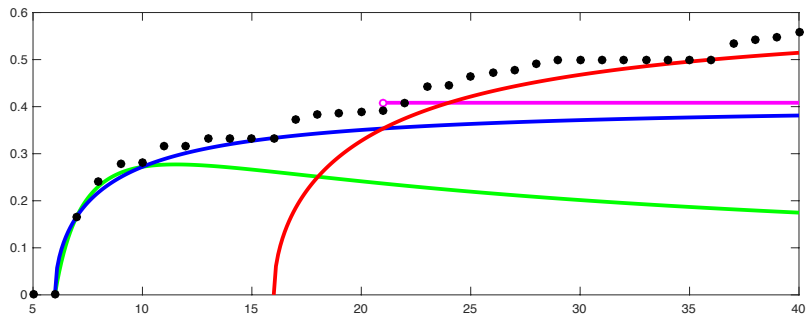


Figure 1: Coherence of best known packings in \mathbb{R}^6 for $5 \leq n \leq 40$ along with best known lower bounds. The Bukh–Cox bound is in green, and the Welch bound is in blue.

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1 The Bukh–Cox Bound

2 The Bukh–Cox Lemma

3 Bukh–Cox Bound via Linear Programming

Setup and Notation

- Let $X = \{x_i\}_{i \in [n]}$ be any sequence in \mathbb{C}^d , and identify X with the $d \times n$ matrix with columns x_i .

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- Denote the set of $\mathbb{C}^{d \times n}$ matrices corresponding to $\frac{n}{d}$ -tight frames by $T(d, n)$.

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- Define $\gamma(d, n) := \max_{Y \in T(d, n)} \|Y^*Y\|_1$.

Bukh–Cox Lemma

Lemma

Let $n = d + k$. Then, every $X \in N(d, n)$ satisfies

$$\mu(X) \geq \frac{n}{\|Y^*Y\|_1 - n} \geq \frac{n}{\gamma(k, n) - n}.$$

Furthermore, X minimizes $\mu(X)$ over $N(d, n)$ if X is equiangular and there exists $Y = \{y_i\}_{i \in [n]} \in T(k, n)$ such that

- (i) Y maximizes $\|Y^*Y\|_1$ over $T(k, n)$,
- (ii) $XY^* = 0$, and
- (iii) $\text{sgn}\langle x_i, x_j \rangle = -\text{sgn}\langle y_i, y_j \rangle$ for $1 \leq i < j \leq n$.

Relation to Welch Bound

Theorem

For all $Y \in T(k, n)$ we have

$$\|Y^*Y\|_1 \leq n + \left[n(n-1) \left(\frac{n^2}{k} - n \right) \right]^{1/2}.$$

Equality is achieved if and only if Y is an equiangular tight frame.

Corollary

Let $n > d$. For all $X \in N(d, n)$,

$$\mu(X) \geq \sqrt{\frac{n-d}{d(n-1)}}.$$

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$\|Y^*Y\|_1$ bound for Bukh–Cox

Theorem

For all $Y \in T(k, n)$ we have

$$\|Y^*Y\|_1 \leq \frac{n^2(1 + (k - 1)\sqrt{1 + k})}{k^2}.$$

Equality is achieved when Y is of the form $Y = [Z|Z|\cdots|Z]$, where $Z \in \mathbb{C}^{k \times k^2}$ is an equiangular tight frame.

Corollary

Let $n > d$. For all $X \in N(d, n)$,

$$\mu(X) \geq \frac{(n - d)^2}{n + (n^2 - nd - n)\sqrt{1 + n - d} - (n - d)^2}.$$

Special Polynomials

To prove the theorem, we will need the following special polynomials:

$$Q_0(x) = 1,$$

$$Q_1(x) = x - \frac{1}{k},$$

$$Q_2(x) = x^2 - \frac{4}{k+2}x + \frac{2}{(k+1)(k+2)}.$$

Proof of Theorem (1/6)

- Without loss of generality assume $y_i \neq 0$ for every $i \in [n]$. First, we normalize the columns of Y , $\{y_i\}_{i \in [n]}$, by defining $z_i := y_i / \|y_i\|_2$.
- The desired bound comes from finding a feasible set of coefficients for the following linear program:

$$\begin{aligned} & \text{minimize} && c_0 \\ & \text{subject to} && f(x) = c_0 Q_0(x) + c_1 Q_1(x) + c_2 Q_2(x), \\ & && 0 \leq c_1 \leq k c_0, c_2 \leq 0, \\ & && f(x) \geq \sqrt{x}, \forall x \in [0, 1]. \end{aligned}$$

Proof of Theorem (2/6)

Suppose we have a feasible (c_0, c_1, c_2) . Then,

$$\begin{aligned}\|Y^*Y\|_1 &= \sum_{i=1}^n \sum_{j=1}^n |\langle z_i, z_j \rangle| \|y_i\|_2 \|y_j\|_2 \\ &\leq \sum_{i=1}^n \sum_{j=1}^n f(|\langle z_i, z_j \rangle|^2) \|y_i\|_2 \|y_j\|_2 \\ &= \sum_{\ell=0}^2 c_\ell \sum_{i=1}^n \sum_{j=1}^n Q_\ell(|\langle z_i, z_j \rangle|^2) \|y_i\|_2 \|y_j\|_2\end{aligned}$$

Proof of Theorem (3/6)

Now we establish a bound for each innermost term $\ell = \{0, 1, 2\}$. Starting with $\ell = 0$, since $Q_0(x) = 1$, we have

$$\sum_{i=1}^n \sum_{j=1}^n Q_0(|\langle z_i, z_j \rangle|^2) \|y_i\|_2 \|y_j\|_2 = \left(\sum_{i=1}^n \|y_i\|_2 \right)^2 := S.$$

For $\ell = 1$, using $Q_1(x) = x - \frac{1}{k}$, we have

$$\sum_{i=1}^n \sum_{j=1}^n Q_1(|\langle z_i, z_j \rangle|^2) \|y_i\|_2 \|y_j\|_2 = \sum_{i=1}^n \sum_{j=1}^n \frac{|\langle y_i, y_j \rangle|^2}{\|y_i\|_2 \|y_j\|_2} - \frac{S}{k}.$$

Proof of Theorem (4/6)

To bound the first term in the $\ell = 1$ case,

$$\begin{aligned} \sum_{i=1}^n \sum_{j=1}^n \frac{|\langle y_i, y_j \rangle|^2}{\|y_i\|_2 \|y_j\|_2} &\leq \left(\sum_{i=1}^n \sum_{j=1}^n \frac{|\langle y_i, y_j \rangle|^2}{\|y_i\|_2^2} \right)^{1/2} \left(\sum_{i=1}^n \sum_{j=1}^n \frac{|\langle y_i, y_j \rangle|^2}{\|y_j\|_2^2} \right)^{1/2} \\ &= \frac{n^2}{k}. \end{aligned}$$

Overall, for $\ell = 1$ we have

$$\sum_{i=1}^n \sum_{j=1}^n Q_1(|\langle z_i, z_j \rangle|^2) \|y_i\|_2 \|y_j\|_2 \leq \frac{1}{k} (n^2 - S).$$

Proof of Theorem (5/6)

For the $\ell = 2$ case, let $\{e_m\}_{m=1}^{d_2}$ be an orthonormal basis for the (finite) vector space spanned by degree-4 projective harmonic polynomials in k variables. There exists a constant $C_{d_2,k}$ such that

$$\begin{aligned} \sum_{i=1}^n \sum_{j=1}^n Q_2(|\langle z_i, z_j \rangle|^2) \|y_i\|_2 \|y_j\|_2 &= C_{d_2,k} \sum_{i=1}^n \sum_{j=1}^n \sum_{m=1}^{d_2} e_m(z_i) \overline{e_m(z_j)} \|y_i\|_2 \|y_j\|_2 \\ &= C_{d_2,k} \sum_{m=1}^{d_2} \left| \sum_{i=1}^n e_m(z_i) \|y_i\|_2 \right|^2 \geq 0. \end{aligned}$$

Multiplying both sides by $c_2 \leq 0$ gives

$$\sum_{i=1}^n \sum_{j=1}^n c_2 Q_2(|\langle z_i, z_j \rangle|^2) \|y_i\|_2 \|y_j\|_2 \leq 0.$$

Proof of Theorem (6/6)

Putting all the bounds together gives

$$\begin{aligned}\|Y^*Y\|_1 &= \sum_{\ell=0}^2 c_\ell \sum_{i=1}^n \sum_{j=1}^n Q_\ell(|\langle z_i, z_j \rangle|^2) \|y_i\|_2 \|y_j\|_2 \\ &\leq c_0 S + c_1 \frac{1}{k} (n^2 - S) \leq c_0 n^2.\end{aligned}$$

Equality is achieved if

- 1 $|\langle y_i, z_j \rangle| = |\langle z_i, y_j \rangle|, \forall i, j,$
- 2 $f(|\langle z_i, z_j \rangle|^2) = |\langle z_i, z_j \rangle|, \forall i, j,$
- 3 $\sum_i \sum_j Q_2(|\langle z_i, z_j \rangle|^2) \|y_i\|_2 \|y_j\|_2 = 0,$
- 4 $\|y_i\|_2 = 1, \forall i,$

which occurs when Y is multiple copies of an ETF of k^2 vectors in \mathbb{C}^k .