# Biangular Gabor Frames and Zauner's Conjecture 

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## Outline

(1) Introduction and Motivation
(2) Proposed Approach
(3) Numerical Experiments

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## Frames

## Definition

$F=\left\{f_{j}\right\}_{j=1}^{n}$ in $\mathbb{C}^{d}$ is a frame if

$$
A\|x\|_{2}^{2} \leq \sum_{j=1}^{n}\left|\left\langle x, f_{j}\right\rangle\right|^{2} \leq B\|x\|_{2}^{2} \quad \forall x \in \mathbb{C}^{d}
$$



Mercedes-Benz

Furthermore, we say $F$ is

- tight if $A=B$ is possible
- unit norm if $\left\|f_{j}\right\|_{2}=1$ for every $j$
- equiangular if $\exists \alpha \geq 0$ such that $\left|\left\langle f_{j}, f_{j^{\prime}}\right\rangle\right|^{2}=\alpha$ whenever $j \neq j^{\prime}$


## Zauner's Conjecture

Equiangular tight frames (ETFs) span optimally packed lines
Many applications:

- Compressed sensing
- Digital fingerprinting
- Multiple description coding

Important question: When do they exist?

## Theorem (Gerzon bound)

There exists an ETF of $n$ vectors in $\mathbb{C}^{d}$ only if $n \leq d^{2}$.

Zauner's Conjecture: For every $d$, there exists an ETF of $d^{2}$ vectors.

## Gabor Frames

## Definition

- Translation operator: $(T v)(j)=v(j-1)$
- Modulation operator: $(M v)(j)=e^{2 \pi i j / d} \cdot v(j)$
- Gabor frame: $G(v):=\left\{M^{\ell} T^{k} v\right\}_{k, \ell=0}^{d-1}$
- Fiducial vector: $v$ such that $G(v)$ is equiangular

Gabor frames are classically used in time-frequency analysis
Zauner's Conjecture (again): There's a fiducial vector in every $\mathbb{C}^{d}(!)$

## Constructive solutions to Zauner

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What about an unconditional proof? We'll need to be non-constructive...

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## Biangular Gabor frames

Big idea: Relax to biangular frames and use intermediate value theorem

## Definition

$G(v)$ is $(\alpha, \beta)$-biangular if
(i) $\left|\left\langle v, T^{k} v\right\rangle\right|^{2}=\alpha$ for $k \in\{1, \cdots, d-1\}$, and
(ii) $\left|\left\langle v, M^{\ell} T^{k} v\right\rangle\right|^{2}=\beta$ for $k \in\{0, \cdots, d-1\}$ and $\ell \in\{1, \cdots, d-1\}$.

## Lemma

If $G(v)$ is an $(\alpha, \beta)$-biangular Gabor frame for $\mathbb{C}^{d}$, then $\alpha+d \beta=\|v\|_{2}^{4}$.

## Examples

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(requires prime $d \geq 5$, "Gabor MUB")

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- $G(v)$ is biangular
$\Longrightarrow G(c v)$ is biangular for $c \in \mathbb{C}^{\times}$


## Real Algebraic Varieties

Define $B_{d}:=\left\{v \in \mathbb{C}^{d}\right.$ for which $G(v)$ is biangular $\}$
Paradoxical observations:

- $B_{d}$ is defined by $\Omega\left(d^{2}\right)$ polynomials over $2 d+2$ real variables
- (Computer) For some $d, B_{d} / \mathbb{C}^{\times}$is one-dimensional


## Lemma

Given $d \in \mathbb{N}$, suppose

- $B_{d}$ is path-connected, and
- there exists a Gabor MUB in $\mathbb{C}^{d}$.

Then there exists a fiducial in $\mathbb{C}^{d}$.


## Proof of concept: $d=2$

It's convenient to define $C_{d}:=\left\{v \in B_{d}: v(0)=1\right\}$

Easy calculation: $C_{2}=$ union of two circles:


$$
\begin{aligned}
& v_{0}=\left[\begin{array}{c}
1 \\
(1+\sqrt{2}) i
\end{array}\right] \\
& v^{\star}=\text { fiducial } \\
& v_{1}=\left[\begin{array}{l}
1 \\
1
\end{array}\right]
\end{aligned}
$$

## Path-Connectedness

Open question: For which $d$ is $B_{d}$ path-connected?

Sufficient condition:

## Lemma

$C_{d}$ is path-connected $\quad \Longrightarrow \quad B_{d}$ is path-connected

Related work:

- Cahill, Mixon, Strawn (2017): FUNTFs are connected
- Needham, Shonkwiler (2018): Symplectic geometry techniques

Remainder of this talk: Numerical evidence of path-connectivity

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## Method

For each $d \in\{2,4,5\}$, do:

- Put $v_{0}:=$ known numerical fiducial (Scott, Grassl 2010)
- Perturb $v_{0}$ and locally minimize $\sum(\text { defining polys })^{2}$ to get $v_{1}$
- For each $j>1$, locally minimize from the perturbation

$$
v_{j}+c \cdot \frac{v_{j}-v_{j-1}}{\left\|v_{j}-v_{j-1}\right\|_{2}}
$$

to get $v_{j+1}$

## Results

$$
d=2
$$


$d=4$



$$
d=5
$$




## Open Problems

- When is $B_{d}$ path-connected?
- We don't really need Gabor MUBs! Instead:

Find $v \in \mathbb{C}^{d}$ such that $G(v)$ is $(\alpha, \beta)$-biangular with $\alpha<\frac{1}{d+1}$

- Can we use $B_{d}$ to find more numerical fiducials?


## Questions?

## Thanks for listening!

